

Implementation of High-Performance MIMO OFDM Receivers

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Contents

- The focus is on MIMO OFDM receivers performing close to the theoretical optimum (with respect to the BER) and their implementation for high data rates
- The course covers
 - Algorithmic aspects (algorithm selection, optimization for implementation)
 - Implementation aspects (ASIC, ASIP, GPP)
 - Mapping aspects
 - (mapping algorithms onto platforms, considering a tradeoff between BER-performance, latency, throughput, and power/energy).



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Frequency Selective MIMO Channels

A frequency selective MIMO channel can described by a time-varying (time t) channel transfer matrix C(t;τ) (i.e. the matrix of channel impulse responses)

$$\mathbf{C}(t;\tau) = \begin{pmatrix} \mathbf{C}_{1,1}(t;\tau) & \cdots & \mathbf{C}_{1,n_{\mathrm{S}}}(t;\tau) \\ \vdots & & \vdots \\ \mathbf{C}_{n_{\mathrm{R}},1}(t;\tau) & \cdots & \mathbf{C}_{n_{\mathrm{R}},n_{\mathrm{S}}}(t;\tau) \end{pmatrix}$$

 $\mathbf{c}_{r,s}(t;\tau)$: time varying complex baseband impulse response between antenna element s of the transmitter and antenna element r of the receiver at time t.

• The receive signal vector (n_R components) is given by $\mathbf{y}(t) = \int \mathbf{C}(t;\tau) \mathbf{s}(t-\tau) d\tau + \mathbf{n}(t)$

with s(t) being the transmit vector (n_s components) and n being an additive distortion (n_R components, e.g. Gaussian noise)

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 T_{svm} : OFDM symbol duration, Δf : subcarrier spacing











MIMO Transmission

Scenario 1

- Channel known to transmitter ("deterministic channel")
- ⇒ Pre-process signal accordingly
 - = Singular-Value-Decomposition (SVD) Architecture
 - Eigen values of **HH**^H : σ_i with i = 1, ..., n_{min} ; n_{min} =min{ n_s , n_R } Singular values: $\lambda_i = \sqrt{\sigma_i}$





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Source: David Tse, Pramod Viswanath, "Fundamentals of Wireless Communication", Cambridge Press 2005





SVD: <u>Singular Value Decomposition as shown on previous slide</u>

Source: David Tse, Pramod Viswanath, "Fundamentals of Wireless Communication", Cambridge Press 2005



MIMO Transmission

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- Optimum power allocation to data streams:
 "waterfilling" allocation



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Receiver Architectures

- SVD transceiver architectures require the channel to be known at the transmitter
 - In time division duplex (TDD) systems the channel may be determined from the uplink due to the channel reciprocity (although this is not trivial since transmitter and receiver typically differ in up- and downlink)
 - Otherwise, the channel must be measured in the receiver and transmitted back to the transmitter via a return link (this is often not feasible; also the channel estimate may be outdated with time-varying channels)
- When the channel is not known at the transmitter we cannot do channel matching preprocessing and, thus, cannot apply SVD.
- \Rightarrow We then also need a different receiver processing.



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- Q=V yields the SVD architecture
- Q=I_{ns}, yields a direct mapping of the data streams to the transmit antennas

Based on: David Tse, Pramod Viswanath, "Fundamentals of Wireless Communication", Cambridge Press 2005









Non-SVD Receiver Architectures

- In this lecture we will focus on non-SVD receiver architectures
 - Channel estimation will not be discussed
 - The MIMO demapper separates the symbol streams, detects the symbols and demapps the bits from the symbols
- We can distinguish fundamental approaches to MIMO demapping
 - with respect to the output
 - providing bits = hard decision output
 - providing bits with reliability information = soft output
 - with respect to the input
 - received signal only
 - received signal and feedback from the decoder
 - bits/symbols
 - bits/symbols with reliability information

iterative receivers









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Non-iterative Demapping

Linear MIMO demappers

Algorithmic options for non-iterative demapping

- Elimination of interference from other data streams: Zero Forcing (ZF) detector
- Minimization of interference plus noise power: Minimum mean square error (MMSE) detector

Maximum Likelihood MIMO demapper

- Hard sphere decoder
- Soft sphere decoder
- Other MIMO demappers (not discussed here)
 - e.g. MCMC





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- appropriate preprocessing).
- If the number of receive antennas is larger than the number of parallel data streams (n_R>n_T) the above receiver inherently exploits receive diversity.



Linear MIMO Demapping: MMSE

Interference nulling focuses on elimination of the interference by the other streams but does not take the noise into account.

A better strategy is minimizing the total distortion z_k(m). For the k-th data stream the effective channel is

$$\mathbf{y}(\mathbf{m}) = \mathbf{H}\mathbf{x}(\mathbf{m}) + \mathbf{w}(\mathbf{m}) = \mathbf{h}_{k}\mathbf{x}_{k}(\mathbf{m}) + \underbrace{\sum_{i \neq k} \mathbf{h}_{i}\mathbf{x}_{i}(\mathbf{m}) + \mathbf{w}(\mathbf{m})}_{\mathbf{z}_{k}(\mathbf{m})}$$



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Linear MIMO Demapping: MMSE

We assume that the channel is known to the receiver and that data streams and noise are uncorrelated

$$\mathbf{R}_{\mathbf{x}\mathbf{x}} = \mathbf{I}_{n_{\mathsf{T}}} \qquad \mathbf{R}_{\mathsf{n}\mathsf{n}} = \mathbf{N}_{\mathsf{0}}\mathbf{I}_{\mathsf{n}_{\mathsf{R}}}$$

The linear minimum mean square error receiver (linear MMSE)

 which equals the MMSE if the x_i have a proper Gaussian distribution – is then given by

$$\hat{\mathbf{x}}(\mathbf{m}) = \mathbf{R}_{\mathbf{x}\mathbf{y}}\mathbf{R}_{\mathbf{y}\mathbf{y}}^{-1}\mathbf{y}(\mathbf{m}) = \mathbf{R}_{\mathbf{x}\mathbf{x}}\mathbf{H}^{\mathsf{H}} \left(\mathbf{N}_{0}\mathbf{I}_{\mathsf{n}_{\mathsf{R}}} + \mathbf{H}\mathbf{R}_{\mathbf{x}\mathbf{x}}\mathbf{H}^{\mathsf{H}} \right)^{-1}\mathbf{y}(\mathbf{m})$$
$$= \left(\mathbf{N}_{0}\mathbf{R}_{\mathbf{x}\mathbf{x}}^{-1} + \mathbf{H}^{\mathsf{H}}\mathbf{H} \right)^{-1}\mathbf{H}^{\mathsf{H}}\mathbf{y}(\mathbf{m})$$

$$\hat{\mathbf{x}}(m) = \mathbf{H}^{H} \left(\mathbf{N}_{0} \mathbf{I}_{n_{R}} + \mathbf{H} \mathbf{H}^{H} \right)^{-1} \mathbf{y}(m) = \left(\mathbf{N}_{0} \mathbf{I}_{n_{T}} + \mathbf{H}^{H} \mathbf{H} \right)^{-1} \mathbf{H}^{H} \mathbf{y}(m)$$

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MIMO Demapping: Maximum Likelihood Detector

The optimum detection scheme with the minimum error probability for equally likely symbols is Maximum Likelihood (ML)

$$\hat{\mathbf{x}} = \underset{\mathsf{all} \mathbf{x}}{\operatorname{arg\,min}} \left\{ \lambda(\mathbf{x}) \right\} = \underset{\mathsf{all} \mathbf{x}}{\operatorname{arg\,min}} \left\{ \left(\mathbf{y} - \mathbf{H} \mathbf{x} \right)^{\mathsf{H}} \left(\mathbf{y} - \mathbf{H} \mathbf{x} \right) \right\} = \underset{\mathsf{all} \mathbf{x}}{\operatorname{arg\,min}} \left\{ \left\| \mathbf{y} - \mathbf{H} \mathbf{x} \right\|^{2} \right\}$$

Efficient implementation as "Hard Sphere Decoder":

$$\mathbf{\hat{x}} = \underset{all \mathbf{x}}{\operatorname{arg\,min}} \{\lambda(\mathbf{x})\} = \underset{all \mathbf{x}}{\operatorname{arg\,min}} \left\{ \left\| \begin{pmatrix} \mathbf{\hat{y}}_{1} & - \sum_{k=1}^{n_{T}} \mathbf{r}_{1k} \mathbf{x}_{k} \\ \vdots \\ \mathbf{\hat{y}}_{n_{T}-1} - \sum_{k=n_{T}-1}^{n_{T}} \mathbf{r}_{(n_{T}-1)k} \mathbf{x}_{k} \\ \mathbf{\hat{y}}_{n_{T}} & - \mathbf{r}_{n_{T}n_{T}} \mathbf{x}_{n_{T}} \end{pmatrix} \right\|^{2} \right\}$$



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The channel matrix H is a n_R×n_T matrix with n_R≥ n_T. Thus, we can triangularize this matrix using the QR decomposition: H=QR, where Q is a n_R×n_T matrix with orthonormal columns and R is an upper triangular n_T×n_T matrix. Since Q has orthonormal columns

 $\boldsymbol{Q}^{\mathsf{H}}\boldsymbol{Q} = \boldsymbol{I}_{\mathsf{n}_{\mathsf{T}}\mathsf{x}\mathsf{n}_{\mathsf{T}}}$

We now expand Q so that it becomes a square matrix, i.e. we append n_R-n_T orthonormal columns (for n_R=n_T no more columns are added)

 $\overline{\mathbf{Q}} = \begin{bmatrix} \mathbf{Q} & \mathbf{Q'} \end{bmatrix}$

Since the new matrix has orthonormal columns and is a square matrix it is unitary ! Also, we expand R by n_R-n_T rows of zeroes so that QR=QR

$$\overline{\mathbf{R}} = \begin{bmatrix} \mathbf{R} \\ \mathbf{0} \end{bmatrix} \implies \overline{\mathbf{Q}} \overline{\mathbf{R}} = \begin{bmatrix} \mathbf{Q} & \mathbf{Q'} \end{bmatrix} \begin{bmatrix} \mathbf{R} \\ \mathbf{0} \end{bmatrix} = \mathbf{Q} \mathbf{R}$$

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- The metric now reads $\lambda(\mathbf{x}) = \|\mathbf{y} - \mathbf{H}\mathbf{x}\|^{2} = \|\mathbf{y} - \mathbf{Q}\mathbf{R}\mathbf{x}\|^{2} = \|\mathbf{y} - \overline{\mathbf{Q}}\mathbf{R}\mathbf{x}\|^{2}$ $= (\mathbf{y} - \overline{\mathbf{Q}}\mathbf{R}\mathbf{x})^{H}(\mathbf{y} - \overline{\mathbf{Q}}\mathbf{R}\mathbf{x})$
- Since $\overline{\mathbf{Q}}$ is unitary $\overline{\mathbf{Q}}\overline{\mathbf{Q}}^{\mathsf{H}} = \overline{\mathbf{Q}}^{\mathsf{H}}\overline{\mathbf{Q}} = \mathbf{I}_{\mathsf{n}_{\mathsf{R}}\mathsf{x}\mathsf{n}_{\mathsf{R}}}$ $\lambda(\mathbf{x}) = (\mathbf{y} - \overline{\mathbf{Q}}\overline{\mathbf{R}}\mathbf{x})^{\mathsf{H}}\overline{\mathbf{Q}}\overline{\mathbf{Q}}^{\mathsf{H}}(\mathbf{y} - \overline{\mathbf{Q}}\overline{\mathbf{R}}\mathbf{x})$ $= (\overline{\mathbf{Q}}^{\mathsf{H}}\mathbf{y} - \overline{\mathbf{R}}\mathbf{x})^{\mathsf{H}}(\overline{\mathbf{Q}}^{\mathsf{H}}\mathbf{y} - \overline{\mathbf{R}}\mathbf{x}) = \|\overline{\mathbf{Q}}^{\mathsf{H}}\mathbf{y} - \overline{\mathbf{R}}\mathbf{x}\|^{2}$
- Replacing Q and R by their definitions

$$\lambda(\mathbf{x}) = \left\| \begin{bmatrix} \mathbf{Q} & \mathbf{Q}' \end{bmatrix}^{\mathsf{H}} \mathbf{y} - \begin{bmatrix} \mathbf{R} \\ \mathbf{0} \end{bmatrix} \mathbf{x} \right\|^{2} = \left\| \begin{bmatrix} \mathbf{Q}^{\mathsf{H}} \mathbf{y} \\ \mathbf{Q}'^{\mathsf{H}} \mathbf{y} \end{bmatrix}^{2} - \begin{pmatrix} \mathbf{R} \mathbf{x} \\ \mathbf{0} \end{pmatrix} \right\|^{2} = \left\| \begin{bmatrix} \mathbf{Q}^{\mathsf{H}} \mathbf{y} - \mathbf{R} \mathbf{x} \\ \mathbf{Q}'^{\mathsf{H}} \mathbf{y} \end{bmatrix} \right\|^{2}$$
$$\lambda(\mathbf{x}) = \left\| \begin{bmatrix} \mathbf{Q}^{\mathsf{H}} \mathbf{y} - \mathbf{R} \mathbf{x} \\ \mathbf{Q}'^{\mathsf{H}} \mathbf{y} \end{bmatrix}^{2} + \left\| \begin{bmatrix} \mathbf{Q}'^{\mathsf{H}} \mathbf{y} \\ \mathbf{Q}'^{\mathsf{H}} \mathbf{y} \end{bmatrix} \right\|^{2}$$

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The task of the ML-receiver is finding the symbol vector x which yields the minimum distance

$$\hat{\mathbf{x}} = \underset{\text{all }\mathbf{x}}{\arg\min} \{\lambda(\mathbf{x})\} \qquad \lambda(\mathbf{x}) = \left\| \left(\mathbf{Q}^{\mathsf{H}} \mathbf{y} - \mathbf{R} \mathbf{x} \right) \right\|^{2} + \left\| \left(\mathbf{Q}^{\mathsf{H}} \mathbf{y} \right) \right\|^{2}$$

Since the right hand term in λ is constant for a given received signal vector we can minimize instead

$$\lambda(\mathbf{x}) = \left\| \left(\mathbf{Q}^{\mathsf{H}} \mathbf{y} - \mathbf{R} \mathbf{x} \right) \right\|^{2} = \left\| \left(\hat{\mathbf{y}} - \mathbf{R} \mathbf{x} \right) \right\|^{2}$$

where

$$\hat{\boldsymbol{y}} = \boldsymbol{Q}^{\mathsf{H}} \boldsymbol{y} = \begin{pmatrix} \sum_{k=1}^{n_{\mathsf{R}}} q_{k1}^{*} \boldsymbol{y}_{k} \\ \vdots \\ \sum_{k=1}^{n_{\mathsf{R}}} q_{kn_{\mathsf{T}}}^{*} \boldsymbol{y}_{k} \end{pmatrix}$$

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The task of the ML-receiver, thus, is finding the symbol vector x which yields the minimum distance

$$\hat{\mathbf{x}} = \underset{all \mathbf{x}}{\operatorname{arg\,min}} \{\lambda(\mathbf{x})\} = \underset{all \mathbf{x}}{\operatorname{arg\,min}} \left\{ \left\| \begin{bmatrix} \hat{\mathbf{y}}_{1} & - \sum_{k=1}^{n_{T}} \mathbf{r}_{1k} \mathbf{x}_{k} \\ \vdots \\ \hat{\mathbf{y}}_{n_{T}-1} - \sum_{k=n_{T}-1}^{n_{T}} \mathbf{r}_{(n_{T}-1)k} \mathbf{x}_{k} \\ \hat{\mathbf{y}}_{n_{T}} & - \mathbf{r}_{n_{T}n_{T}} \mathbf{x}_{n_{T}} \end{bmatrix} \right\|^{2}$$

Where we have made use of the fact that R is an upper triangular matrix

$$\mathbf{R} = \begin{pmatrix} \mathbf{r}_{11} & \cdots & \cdots & \mathbf{r}_{1n_{T}} \\ 0 & \ddots & & \vdots \\ \vdots & \ddots & \ddots & \vdots \\ 0 & \cdots & 0 & \mathbf{r}_{n_{T}n_{T}} \end{pmatrix}$$

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$$\hat{\mathbf{x}} = \underset{all \mathbf{x}}{\operatorname{arg\,min}} \{\lambda(\mathbf{x})\} = \underset{all \mathbf{x}}{\operatorname{arg\,min}} \left\{ \left\| \begin{pmatrix} \hat{\mathbf{y}}_{1} & - \sum_{k=1}^{n_{T}} \mathbf{r}_{1k} \mathbf{x}_{k} \\ \vdots \\ \hat{\mathbf{y}}_{n_{T}-1} - \sum_{k=n_{T}-1}^{n_{T}} \mathbf{r}_{(n_{T}-1)k} \mathbf{x}_{k} \\ \hat{\mathbf{y}}_{n_{T}} & - \mathbf{r}_{n_{T}n_{T}} \mathbf{x}_{n_{T}} \end{pmatrix} \right\|^{2}$$

The squared norm of a vector equals the sum of the squared magnitudes of its components

$$\left\|\mathbf{z}\right\|^{2} = \left\| \begin{pmatrix} z_{1} \\ \vdots \\ z_{n} \end{pmatrix} \right\|^{2} = \sum_{k=1}^{n} \left| z_{k} \right|^{2} \qquad \Rightarrow \qquad \lambda(\mathbf{x}) = \sum_{m=1}^{n_{T}} \left| \hat{\mathbf{y}}_{m} - \sum_{k=m}^{n_{T}} \mathbf{r}_{mk} \mathbf{x}_{k} \right|^{2}$$

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The task of the ML-receiver, thus, is finding the symbol vector x which yields the minimum distance

$$\hat{\mathbf{x}} = \underset{\text{all }\mathbf{x}}{\operatorname{arg min}} \left\{ \lambda(\mathbf{x}) \right\} = \underset{\text{all }\mathbf{x}}{\operatorname{arg min}} \left\{ \sum_{m=1}^{n_{T}} \left| \hat{\mathbf{y}}_{m} - \sum_{k=m}^{n_{T}} r_{mk} \mathbf{x}_{k} \right|^{2} \right\}$$

We note two important facts

- The last term of the summation (m=n_T) only depends on symbol x_{n_T} and with decreasing m the other symbols in the vector x step-by-step become terms of the summation
- The contribution of each component of the vector to the total sum is non-negative, i.e. ≥0 ! (since we add the squared magnitude of this component)
- ⇒ This is the basis of the sphere decoding algorithm



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The task of the ML-receiver, thus, is finding the symbol vector x which yields the minimum distance

$$\hat{\mathbf{x}} = \underset{\text{all }\mathbf{x}}{\operatorname{arg min}} \{ \lambda(\mathbf{x}) \} = \underset{\text{all }\mathbf{x}}{\operatorname{arg min}} \left\{ \sum_{m=1}^{n_{T}} \left| \hat{\mathbf{y}}_{m} - \sum_{k=m}^{n_{T}} r_{mk} \mathbf{x}_{k} \right|^{2} \right\}$$

- 1. We evaluate the summation terms running with index m in reverse order, i.e. from n_T to 1, for one possible data vector x_1
- 2. This yields a metric $\lambda(x_1)$ which we use as reference $\lambda(x_{ref})$
- 3. We do the same for a second possible data vector x₂. Whenever the intermediate metric value is larger than λ(x_{ref}), we can abort the calculation of λ(x₂). If a smaller λ results, we use this as the new reference λ(x_{ref}).
 4. Finally, the last λ(x_{ref}) is the metric of the ML estimate x̂ = arg{λ(x_{ref})}
- A clever search order significantly reduces the amount of calculations (e.g. Schnorr-Euchner enumeration)

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Sphere Decoder - Comments

- With the method shown on the previous slide we usually do not have to calculate the metrics for all branches/leafs of the tree. The reduction of the computational effort can be very significant
- The savings in computation depend on channel matrix and noise. Therefore, the actual computation time cannot be predicted. This can cause problems with latency and throughput.
- There are various approaches reducing latency and increasing throughput at the cost of slight performance losses (which will not be discussed here).





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Soft Output • Log Likelihood Ratio (LLR) $L(c_{n,b}|\mathbf{y}) = ln \frac{P(c_{n,b} = 0|\mathbf{y}, \mathbf{H})}{P(c_{n,b} = 1|\mathbf{y}, \mathbf{H})} = ln \frac{\sum_{\mathbf{x} \in S_{n,b}^{0}} e^{-\frac{\|\mathbf{y} - \mathbf{H}\mathbf{x}\|^{2}}{N_{0}}}}{\sum_{\mathbf{x} \in S_{n,b}^{1}} e^{-\frac{\|\mathbf{y} - \mathbf{H}\mathbf{x}\|^{2}}{N_{0}}}}$

The LLR computation is very complex and has to be performed for each bit. Therefore, commonly the max-log approximation is used

$$\mathbf{L}(\mathbf{c}_{n,b}|\mathbf{y}) \approx \frac{1}{N_0} \left(\min_{\mathbf{x} \in S_{n,b}^1} \|\mathbf{y} - \mathbf{H}\mathbf{x}\|^2 - \min_{\mathbf{x} \in S_{n,b}^0} \|\mathbf{y} - \mathbf{H}\mathbf{x}\|^2 \right)$$

 $S_{n,b}^{0/1}$ =set of symbols for which $c_{n,b} = 0/1$



Soft Output: Linear MIMO Demapping

Max-log approximation

$$L(c_{n,b}|\mathbf{y}) \approx \frac{1}{N_0} \left(\min_{\mathbf{x} \in S_{n,b}^1} \|\mathbf{y} - \mathbf{H}\mathbf{x}\|^2 - \min_{\mathbf{x} \in S_{n,b}^0} \|\mathbf{y} - \mathbf{H}\mathbf{x}\|^2 \right)$$

For ZF the LLR can be approximated by

$$L(c_{n,b}|\hat{x}_{n}) \approx \frac{1}{\sigma_{\hat{w},n}^{2}} \left(\min_{\mathbf{x}_{n} \in S_{n,b}^{1}} \|\hat{x}_{n} - x_{n}\|^{2} - \min_{\mathbf{x}_{n} \in S_{n,b}^{0}} \|\hat{x}_{n} - x_{n}\|^{2} \right)$$

• For MMSE the LLR can be approximated by $L(c_{n,b}|\hat{x}_{n}) \approx \frac{1}{MSE_{MMSE,n}} \left(\min_{\mathbf{x}_{n} \in S_{n,b}^{1}} \|\hat{x}_{n} - x_{n}\|^{2} - \min_{\mathbf{x}_{n} \in S_{n,b}^{0}} \|\hat{x}_{n} - x_{n}\|^{2} \right)$

 $S_{n,b}^{0/1}$ =set of symbols for which $c_{n,b} = 0/1$

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Soft Output: Sphere Decoder

Max-log approximation

$$L(\mathbf{c}_{n,b}|\mathbf{y}) \approx \frac{1}{N_0} \left(\min_{\mathbf{x} \in S_{n,b}^1} \|\mathbf{y} - \mathbf{H}\mathbf{x}\|^2 - \min_{\mathbf{x} \in S_{n,b}^0} \|\mathbf{y} - \mathbf{H}\mathbf{x}\|^2 \right)$$

This requires minimum searches over all symbols x where c_{n,b} =0 and where c_{n,b} =1. The two searches may be combined into a single tree search (STS).

 $S_{n,b}^{0/1}$ =set of symbols for which $c_{n,b} = 0/1$

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Hard Decision Feedback: Successive Interference Cancellation (SIC)

Not really iterative!

We may use ZFE or MMSE, detect the data stream with the largest SNR (if the data stream is coded, the decision error probability can be very low). Then, the contribution of this data stream is subtracted from the received signal and the process is repeated until all data streams are detected.



Source: David Tse, Pramod Viswanath, "Fundamentals of Wireless Communication", Cambridge Press 2005





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Successive Interference Cancellation (SIC)

Not really iterative!

After subtraction (assuming that we have detected correctly, i.e. x₁=x₁), we can use a new receive signal model for the de-correlation of the second data stream:

$$\mathbf{y}_{1}(\mathbf{m}) = \mathbf{y}(\mathbf{m}) - \mathbf{H} \begin{pmatrix} \hat{\mathbf{x}}_{1}(\mathbf{m}) \\ 0 \\ \vdots \\ 0 \end{pmatrix} = \mathbf{H} \begin{pmatrix} 0 \\ \mathbf{x}_{2}(\mathbf{m}) \\ \vdots \\ \mathbf{x}_{n_{T}}(\mathbf{m}) \end{pmatrix} + \mathbf{w}(\mathbf{m}) = \mathbf{H}_{1}\mathbf{x}_{1}(\mathbf{m}) + \mathbf{w}(\mathbf{m})$$

with

$$\mathbf{x}_{1}(\mathbf{m}) = \begin{pmatrix} \mathbf{x}_{2}(\mathbf{m}) \\ \vdots \\ \mathbf{x}_{n_{T}}(\mathbf{m}) \end{pmatrix} \quad ; \quad \mathbf{H}_{1} = \begin{pmatrix} \mathbf{h}_{2} & \cdots & \mathbf{h}_{n_{T}} \end{pmatrix}$$

where h_k are the column vectors of H=($h_1 h_2 ... h_{n_T}$)



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Considering Soft Feedback in Sphere Decoding

Max-log approximation

$$L(\mathbf{c}_{n,b}|\mathbf{y}) \approx \min_{\mathbf{x}\in S_{n,b}^{1}} \left\{ \frac{\|\mathbf{y} - \mathbf{H}\mathbf{x}\|^{2}}{N_{0}} - \sum_{\nu=1}^{n_{T}} \sum_{\beta=1}^{B} \ln P(\mathbf{c}_{\nu,\beta}) \right\} - \min_{\mathbf{x}\in S_{n,b}^{0}} \left\{ \frac{\|\mathbf{y} - \mathbf{H}\mathbf{x}\|^{2}}{N_{0}} - \sum_{\nu=1}^{n_{T}} \sum_{\beta=1}^{B} \ln P(\mathbf{c}_{\nu,\beta}) \right\}$$

This requires minimum searches over all symbols x where c_{n,l} =0 and where c_{n,l} =1. The two searches may be combined into a single tree search (STS).

 $S_{n,b}^{0/1}$ =set of symbols for which $c_{n,b} = 0/1$

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Studer, C., Fateh, S., and Seethaler, D.: ASIC Implementation of Soft-Input Soft-Output MIMO Detection Using MMSE Parallel Interference Cancellation



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 $S_{n,b}^{0/1}$ =set of symbols for which $c_{n,b} = 0/1$





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General Implementation Aspects

- Tree Search, occurs in
 - Sphere decoding
 - LLR computation

Matrix, vector algebra

- Scalar product
- Matrix-vector and matrix-matrix multiplication (incl. Hermitian)
- Matrix inversion
 - Direct inversion only reasonable for dim 2
 - Various alternatives via, e.g.,
 LU-decomposition, QR-decomposition
- QR decomposition
 - using Gram-Schmidt of Givens Rotation

What are the main processing functions ?

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Source: Eberli, S., "Application-Specific Processor for MIMO-OFDM Software-Defined Radio", Dissertation, ETH Zurich, 2009





Demapper Challenge: Tree Search

Pruning requires sorting (*enumeration*) of constellation points

Enumeration

- No soft-input: geometrical properties of *M*_C allow efficient enumeration
- Soft-input soft-output: geometrical properties destroyed by a priori information *M*_A



Metrics in Constellation Diagram:

- \mathcal{M}_{C} : Euclidean distance to received symbol **y**
- \mathcal{M}_{A} : Based on soft inputs from channel decoder
- \mathcal{M}_{P} : Tree branch metrics \mathcal{M}_{A} + \mathcal{M}_{C}











Agenda

- Introduction
- MIMO Receivers

Implementation

- General implementation aspects
- Ð

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- Cae²sar: A Scalable VLSI Architecture for SISO Depth-First Sphere Decoding
- IteRX: Silicon Implementation of Iterative MIMO Detection and Decoding

ASIC Implementation of soft-in-soft-out sphere detector for high data rates

- Mapping to Application Specific Platforms
- Summary



A 772 Mbit/s 8.81 bit/nJ 90 nm CMOS Soft-Input Soft-Output Sphere Decoder

Filippo Borlenghi, Ernst Martin Witte, Gerd Ascheid, Heinrich Meyr, Andreas Burg



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Telecommunications Circuits Laboratory







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Sphere Decoding

- Tree search-based detection can be solved by depth-first sphere decoding (SD)
- Branch & bound algorithm
- In a single tree search (STS), SD finds:
 - The maximum-a posteriori (MAP) solution, i.e. the path with the minimum total metric
 - The counter hypothesis for each bit



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7 ICE **High-Level Architecture Basic principle: check** \triangleright one node per cycle (ONPC) **Concurrent computation of** the best child and sibling Horizontal **Vertical** step **Control FSM** step 2 3 5 6 7 cycle 1 4 V-step \sum Þ H-step PC chk Pruning-criteria checks 🚬 2.2 2.3 2.4 3.2 1.2 1.3 1.4 3.3 4.2 62 夏季 營 (二〇一二年 七月 六日) ·學

7 ICE **High-Level Architecture Basic principle: check** \triangleright one node per cycle (ONPC) **Concurrent computation of** the best child and sibling Horizontal **Vertical** step **Control FSM** step 2 3 5 6 7 cycle 1 4 V-step 1 \sum Þ H-step PC chk Pruning-criteria checks 🚬 2.2 2.3 2.4 3.2 1.2 1.3 1.4 3.3 4.2 63 夏季 營 (二〇一二年 七月 六日) -壆

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High-Level Architecture

3.1 3.2

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3.3

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 Basic principle: check one node per cycle (ONPC)

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 Concurrent computation of the best child and sibling

cycle	1	2	3	4	5	6	7
V-step	1	1.1			2.1		3.1
H-step		2	1.2	1.3	3	2.2	4
PC chk		1	1.1	☆	2	201	X

2.2

2.3 2.4

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1.2 1.3 (1.4)

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4.2



Channel-based Enumeration

- Vertical step: quantization
- Horizontal step:
 - 1. Column-wise partition of QAM constellation
 - 2. Order within a column follows zig-zag pattern

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3. Select minimum M_C among all columns



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Silicon Implementation



- Technology: 90nm1.0V UMC
- 3 configurable cores (all up to 4x4 antennas):
 - 4 QAM
 - Up to 16 QAM
 - Up to 64 QAM

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Implementation Results

		4 QAM	16 QAM	64 QAM
Area	[mm ²]	0.27	0.54	0.97
	[kGE]	58	114	213
Max. Frequency	[MHz]	330	244	193
Max. Throughput	[Mbit/s]	440	651	772
Max. Area Efficiency	[Mbit/s/kGE]	7.63	5.73	3.62
Max. Energy Efficiency	[bit/nJ]	8.29	7.82	8.81
Gate equivalent (GE)		C		S
= area of a 2-input NAND gate	Area	X	2	x2
	Frequency	-20	6% -2	21%

- Energy penalty for detecting on 64-QAM core
 - 16 QAM is up by 24% (w.r.t. 16-QAM core)
 - 4 QAM is up by 119% (w.r.t. 4-QAM core)

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Peak Performance Comparison

		This work	Studer 2011	Liao 2010	Shab. 2009
N. of Antennas		≤ 4x4	≤ 4x4	≤ 8x8	4x4
Modulation Order		≤ 64	≤ 64	≤ 64	64
Iterative Det.		YES	YES	NO (soft)	NO (hard)
CMOS Tech.	[nm]	90	90	130	130
Supply Voltage	[V]	1.0	1.2	1.3	1.3
Area	[kGE]	213 ^a	410	350 ^a	114 ^a
Max. Throughput	[Mbit/s]	772	757	624 ^b	946 ^b
Max. Area Eff.	[Mbit/s/kGE]	3.62	1.85	1.78 ^b	8.30 ^b
Max. Energy Eff.	[bit/nJ]	8.81	4.00	18.11 ^b	12.21 ^b

^a QRD required but not included since not executed at symbol rate

 $^{\rm b}$ Scaled to 90nm 1.0V technology by f ~ S and P ~ 1/U^2

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SISO Detectors in Quasi-Static Case

- SD can operate at 1% PER at 3dB lower SNR than MMSE PIC, at reduced throughput (5.7 Mbit/s)
- Let's compare at same SNR and target PER:
 - @ 24dB (min for 1% PER for MMSE-PIC @ 6 it.)

	lts.	Throughput [Mbit/s]	Area Eff. [Mbit/s/kGE]	Energy Eff. [bit/nJ]
STS SD	2	96	0.45	1.04
MMSE PIC	6	126	0.31	0.67

@ 25dB (min for 1% PER for MMSE-PIC @ 4 it.)

	lts.	Throughput [Mbit/s]	Area Eff. [Mbit/s/kGE]	Energy Eff. [bit/nJ]
STS SD	1	154	0.73	1.76
MMSE PIC	4	189	0.46	1.00

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Cae²sar Summary

- Cae²sar chip (2010, E.M. Witte and F. Borlenghi RWTH Aachen)
 - First silicon implementation of SISO depth-first sphere decoding
- Features
 - Max-log MAP optimal detection
 - Variable run-time and adjustable complexity-performance tradeoff
 - Up to 4x4 antennas/64-QAM
- Implementation results in 90 nm technology
 - Area: 212 kGE
 - Max. frequency @ 1.0 V:193 MHz
 - Max. (uncoded) throughput: 772 Mbit/s



Witte et al., A Scalable VLSI Architecture for Soft-Input Soft-Output Single Tree-Search Sphere Decoding, TCAS-II, 2010 Borlenghi et al., A 772 Mbit/s 8.81 bit/nJ 90 nm CMOS Soft-Input Soft-Output Sphere Decoder, A-SSCC 2011

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The Channel Decoder: 802.11n LDPC

- Eagle chip (2009, C. Roth ETH Zurich)
- Features
 - Layered offset-min-sum (OMS) decoding
 - 802.11n compliant
 - Code rates: 1/2, 2/3, 3/4, 5/6
 - Code-word lengths:
 648, 1296, 1944
 - Optimized for low power
 - Early termination support
- Implementation results in 90 nm technology
 - Area: 398 kGates
 - Max. frequency @ 1.0 V: 346 MHz



Max. (coded) throughput: 679 Mbit/s (code rate 5/6, 10 its.)

Roth et al., A 15.8 pJ/bit/iter Quasi-Cyclic LDPC Decoder for IEEE 802.11n in 90 nm CMOS, A-SSCC 2010

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Building an Iterative Detection and Decoding System

- Block-wise processing
 - Detector and decoder exchange complete frames
 - Mutually exclusive access to the current frame

Interleaved ping-pong scheduling

 Detector and decoder process concurrently different frames

\rightarrow Both always busy!





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7 ICE **Building an Iterative Detection and Decoding System** Typically it looks more like this: DETECTOR DETECTOR . . . DECODER DECODER \rightarrow Multiple SD cores needed! **Open questions** How many cores? As many as possible...

- How to distribute the input data?
- How to schedule the execution?
- How to collect the output results?











Latch-Based LLR Memory Architecture

- Latch-based memories introduced in original LDPC design
- Why?
 - Area similar to macro-based (for this size and technology)
 - Easy to customize (multiple addresses within first bank)
 - Low-power



Meinerzhagen et al., Towards Generic Low-Power Area-Efficient Standard Cell Based Memory Architectures

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Implementation Results

- Core area: 2.78 mm² (1.58 MGE) in low-power 65 nm technology
 - Detector (5 SDs): 872 kGE (55%), 140 to 145 kGE / SD
 - Decoder: 447 kGE (28%)
 - LLR memory: 210 kGE (13%)
- Max. frequencies @ 1.2 V
 - Detector: 135 MHz
 - Decoder: 299 MHz
- Supports
 - {2x2, 3x3, 4x4} antennas
 - 4, 16, 64 QAM
 - All 802.11n LDPC codes
- Max. throughput > 1 Gbit/s



Borlenghi et al., A 2.78 mm² 65 nm CMOS Gigabit MIMO Iterative Detection and Decoding Receiver, ESSCIRC 2012

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In cooperation with:

Prof. Andreas Burg





Integrated Systems Laboratory & Microelectronics Design Center



Eidgenössische Technische Hochschule Zürich Swiss Federal Institute of Technology Zurich



Agenda

- Introduction
- MIMO Receivers
- Implementation

Mapping to Application Specific Platforms

Motivation

Implementation using MPSoC platforms

- Nucleus Methodology
- Case Study : MIMO OFDM Transceiver
- Tool Flow

Summary













Agenda

- Introduction
- MIMO Receivers
- Implementation

Mapping to Application Specific Platforms

- Motivation
- Nucleus Methodology
 - Case Study : MIMO OFDM Transceiver
 - Tool Flow

Summary

 \rightarrow

Implementation using MPSoC platforms: we need a mapping methodology !

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Agenda

- Introduction
- MIMO Receivers
- Implementation

Mapping to Application Specific Platforms

- Motivation
- Nucleus Methodology
- Case Study : MIMO OFDM Transceiver
 - Tool Flow

Implementation using MPSoC platforms: Does the proposed methodology work ?

Summary

 $\left(\rightarrow \right)$





Case Study: MIMO OFDM Transceiver



- **Application analysis Nuclei identification**
- Mapping on a ST P2012 platform
- Mapping on a TI-based platform
- Using a preprocessing ASIP







- Outer Modem
 - Channel (De-)coding
 - (De-)Interleaving
- Inner Modem (RX)
 - RX OFDM Processing
 - Channel Estimation
 - Spatial Equalizing: Mitigate channel impact on payload
 - Soft Demapping: Calculate soft bits (LLRs) BPSK, 4QAM, 16QAM, 64QAM



Nuclei Identification: Kernel Identification



Analyze different algorithmic choices within RX blocks

- Identify computational kernels
 - Recurring tasks
 - Operate on data with certain alignment

Build application as composition of kernels




Nuclei Identification: Kernel Overview



- Wide variety of algorithms is covered for key transceiver functions
 - Channel Estimation
 - Spatial Equalization
 - Channel Coding



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What Describes a Nucleus Library Element?

Flavors and configuration parameters influence performance properties

- Processing properties
 - Latency
 - Throughput
 - Energy efficiency
- Implementation properties
 - Area
 - Memory
- Algorithmic properties
 Bit error rate

V. Ramakrishnan et al.: "Efficient Implementations From Libraries: Analyzing the Influence of Configuration Parameters on Key Performance Properties", IEEE PIMRC Conference 2009

			Configuration Parameters				
Case	PE		Input width (bits)	Twiddle width (bits)	Internal scaling	Rounding / Truncation	Cycles per frame
1	C64x	R-4	16	16	yes	R	6526
2		R-4	16	16	no	т	6262
3	C62x	R-2	16	16	no	т	22388
4		R-4	32	16	no	R	8327
5	C64x	R-4	32	32	no	R	11895
6		R-4	32	32	yes	R	11895
7		R-2	16	16	yes	R	7364
8		R-2	16	16	yes	R	12453
9		R-4	16	16	yes	R	3450
10	Virtey-5	R-2	16	16	yes	R	3199
11	THEOR V	R-4	16	16	yes	Т	
12		R-4	16	16	no	Т	
13		R-4	16 (12)	16	yes	R	3450
14		R-4	32	16	yes	R	
15		R-4	32	32	yes	R	-



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Case Study: MIMO OFDM Transceiver

- Application analysis Nuclei identification
- Mapping on a ST P2012 platform
 - IEEE 802.11n based SDR application (EU funded project "2PARMA")
- Mapping on a TI-based platform
- Using a preprocessing ASIP



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- Hardware synchronizer for inter-core signaling
- Interface for hardware accelerators (ASICs)



• Outer Modem: Bitwise data

- Channel Coding
- Interleaving
- Inner Modem: Complex baseband data
 - OFDM Processing
 - Channel Estimation: Using upfront block preamble
 - Spatial Equalizing: Mitigate channel impact on payload
 - Soft Demapping: Calculate LLRs BPSK, 4QAM, 16QAM, 64QAM





Modified Gram Schmidt QRD (MGS-QRD)

- Problem 1: Fixed point range
 - Dynamic Scaling (DS): Keeps values in range
- Problem 2: High SNR
 - □ High data rates (4x4 64QAM) require SNR ≈ 30dB
 - Noise spectral density becomes low

$$\overline{\mathbf{H}} = \begin{pmatrix} \hat{\mathbf{H}} \\ N_0 \mathbf{I} \end{pmatrix} = \begin{pmatrix} \mathbf{Q}_{\mathbf{a}} \\ \mathbf{Q}_{\mathbf{b}} \end{pmatrix} \mathbf{R} \qquad \mathbf{G} = \frac{1}{N_0} \mathbf{Q}_b \mathbf{Q}_a^{\ H} \qquad \widehat{\mathbf{x}} = \mathbf{G} \mathbf{y}$$



Flavor of Nucleus for Numerical Stability – High SNR

A key issue: fixed point

Solution

- Scale up lower part of regularized channel matrix
- Maintain vector norms on the fly (compare ETH alg.)
- Scale down lower part of column vectors for projections

$$\overline{\mathbf{H}}_n = \begin{pmatrix} \hat{\mathbf{H}} \\ \mathbf{I} \end{pmatrix}$$

Givens Rotation QRD (GR-QRD)

- Orthogonalize by vector rotations
- CORDIC algorithm implemented
- Advantages
 - Inherent numerical stability
 - Requires less operations than MGS-QRD



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Small gap between floating point reference and fixed point implementation



- Modified Gram Schmidt QRD (MGS-QRD)
 - Less conditional jumps, less norm shifting in DS
- Equalizer Matrix calculation $G = \mathbf{R}^{-1} \mathbf{Q}_a^H$
 - Use triangular structure of R

. . .

- Simplified 64QAM soft demapping [1]
 - Approximate piecewise linear functions
 - Calculate bit value from previous one from same symbol
 - [1] Simplified Soft-Output Demapper for Binary Interleaved COFDM with Application to HIPERLAN/2, *Filippo Tosato*, HP Laboratories Bristol

Application Implementation: Kernel Overview

	System	2x2		4x4		
	Operation	cycles	time	cycles	time	
			(μs)		(μs)	
	Matrix/Vector Operations					
1	mat-mat add	13	0.022	23	0.038	
2	mat hermitian	26	0.043	90	0.150	
3	mat-rscal mul	26	0.043	45	0.075	
4	mat-vec mul	44	0.073	70	0.117	
5	mat-mat mul	102	0.170	301	0.502	
6	mat-rmat mul	74	0.123	205	0.342	
7	mat-tmat mul	52	0.087	239	0.398	
8	mat-mat mul 8 vm 2^2	218	0.363	503	0.838	
9	mat inv	385	0.642	1,328	2.213	
10	tri mat inv	43	0.072	278	0.463	
11	reg-mgs-qrd	447	0.745	991	1.652	
12	reg-mgs-ds-qrd	703	1.172	1,368	2.280	
13	back subst.	954	1.590	2,106	3.510	
14	back subst. slicing	1,170	1.950	2,538	4.230	
OFDM slot wise operations						
15	bpsk soft demap	329	0.548	658	1.097	
16	4qam soft demap	658	1.097	1,316	2.193	
17	16qam soft demap	857	1.428	1,705	2.842	
18	64qam soft demap	1,559	1.428	3,117	5.195	
19	64qam soft demap appr	1,134	1.890	2,268	3.780	
20	fft	1,774	2.957	3,548	5.913	
21	fft mem realign	2,052	3.420	4,084	6.838	
22	ifft	2,028	3.380	4,056	6.760	
10		•	•	•		

For 2x2 and 4x4 MIMO use case

- Cycles for execution on single STxP70 processor core including VECX unit
- Corresponding time for 600MHz clock frequency

 In the range of ...
 IEEE 802.11n real time (4µs per OFDM slot)

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7 ICE **Occupation Graph** Implementation on P2012 platform using 8 cores Minimum latency for 21 or more OFDM slots of data payload Latency = 8.2μ s, IEEE 802.11n allows 16 μ s (including MAC layer) Real time 64QAM demapping possible \mathcal{D} PE 1 Р Mod 1 PE 2 Mod 2 PE 3 **Spatial Equalizing** Preprocessing PP/EQ 1 PE 4 PP/EQ 2 PE 5 PP/EQ 3 \Box PE 6 PP/EQ 4 PE 7 Symbol Demapping SD 1 PE 8 SD 2 122 16 32 48 80 t[us] 0 64 96







Execution Time 2x2

Component	Cycles	Time (us) @850MHz	
Preprocessing Steps			
OFDM Demodulation	1,296	1.525	
Subcarrier Demapping	1,220	1.435	
CP remove	360	0.424	
Channel Estimation (MIMO, LS)	3,272	3.849	
Detection Preproc. (MIMO 2x2, MMSE, DS)	39,499	46.469	
Sum	45,647	53.702	
Detection Steps			
OFDM Demodulation	648	0.762	
Subcarrier Demapping	610	0.359	
CP remove	180	0.718	
Equalizing - MIMO	984	1.158	
Soft Demapping 4QAM (per OFDM sym.)	364	0.428	
Soft Demapping 16QAM (per OFDM sym.)	618	0.727	
Soft Demapping 64QAM (per OFDM sym.)	1,062	1.249	
Sum	3,484	4.099>	







Alternative 1: ASIP

Key architectural features

- Relatively large register file (64 register) to minimize memory accesses during computations.
- Short 5-stage pipeline to avoid long stall cycles or NOP operations.
- Native support for complex data types and all common arithmetic operations,
- 2-way SIMD for all instructions operating on complex values including random register access.
- Special instructions for computation of reciprocal, inverse square root and loglikelihood-ratio (LLR) computation of Gray-coded modulations
- Design time flexible data path, e.g. high precision fix point format 8.16 bit (=24 bit) for real/imaginary data.
- Two external simple handshake protocol interfaces with separate read and write ports.
- Small data memory (1k words with same data width as data path) to store look-up tables and intermediate results.
- Instructions for simple control mechanisms, e.g. loops and if-else statements.
- Advanced synchronization features based on efficient IRQ In/Out interfaces



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Initial Synthesis Results

Initial synthesis result for 24bit data path (registers, multipliers, etc.)*

Total Core Area	Area Registers	Program Memory	Frequency	
~205 kGE	~88 kGE	~28 kGE (1024 instructions)	400 MHz	

Execution times

Flavor (4x4 antenna configuration)	Cycles (a	pprox.)	Exec. Time in ns
QRD-UDS (augmented matrix) per decomposition		151	377.5
Computation of equalizer matrix G $(G=Q_bQ_a^H)$	19		47.5
SINR computation per symbol vector	36		90.0
Spatial equalization $(x=Gy=(Q_bQ_a^H)y)$ per symbol vector	23		57.5
LLR computation per symbol vector	4QAM	13	32.5
	16QAM	19	47.5
	64QAM	25	62.5

Speed up (compared to P2012) ≈ 1 order of magnitude

* Synopsys Design Compiler E2010.12-SP2 Faraday 90nm Library, Standard Performance, Typical case library & conditions, 1.00V, 25°C

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Alternative Solution 2: CGRA

- Pipelined Coarse Grained Reconfigurable Array (CGRA)
 - Processing Element (PE)
 - Complex multiplier
 - Complex ALU
 - Shifter
 - Local register file
 - PE 2x2
 - MESH connected
 - Broadcast inputs
 - PE Chain
 - Results from PE 2x2
 - Center Alpha Unit
 - Special *alpha* parameter in matrix inversion





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Results: Area and Throughput

Synthesis results*

Component	Area / kGE (Gate Equivalents)
Complex multiplier	13,0
Complex adder	1,1
Processing Element	18,6
CGRA	393
Configuration Memory	43

Processing performance

- Input: H, y
- Output: soft $\hat{\mathbf{x}} = (\mathbf{H}^{\mathbf{H}}\mathbf{H} + N_0\mathbf{I})^{-1}\mathbf{H}^{\mathbf{H}}\mathbf{y}$
- Throughput: Speed up > 2 orders of magnitude (150 Msym/s @ 588 MHz [^]≈ 30 ns/sym)
- * Synopsys Design Compiler Ultra 2010.12-SP2 Faraday 65nm technology, Standard Performance, typical case library & conditions, 1.00V, 25°C

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Tool Flow: Overview



* For further details see the presentation by Rainer Leupers Programming Heterogeneous MPSoC Platforms: *The MAPS Approach*

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Tool Flow: Overview (2)

Screenshot of Graphical User Interface:



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Summary

- A variety of algorithms is available for MIMO receivers
- Computationally more complex, near-optimal algorithms provide superior performance. Using them in iterative receiver architectures yields further performance gains
- Appropriate ASIC architectures enable the energy efficient implementation of iterative receivers for highest data rates
- Advanced MPSoC platforms enable the flexible mapping of different receivers, the more application specific the platform the more efficient
- A limited set of processing functions ("nuclei") can support efficient implementation of a large variety of algorithms
- The "nucleus" methodology represents a systematic design approach for flexible and efficient transceiver design, as demonstrated
 - by the mapping of a MIMO OFDM transceiver onto different MPSoC platforms and
 - by a prototype tool set

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