# A fast approach for overcomplete sparse decomposition based on smoothed L0 norm 

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## Short summary:

This paper proposes a fast algorithm for overcomplete sparse decomposition. The algorithm is derived by directly minimizing the L0 norm after smoothing. Hence, the algorithm is named as smoothed L0 (SL0) algorithm. The authors demonstrate that their algorithm is 2-3 orders of magnitude faster than the state-of-the-art interior point solvers with same (or better) accuracy.

## I. Introduction

- To introduce the algorithm the authors have used the context of source component analysis (SCA). SCA is a method to achieve separation of sparse sources.
- Suppose that $m$ source signals are recorded by a set of $n$ sensors each of which records a combination of all sources. In linear instantaneous (noiseless) model, it is assumed that $x(t)=A s(t)$ in which $s(t)$ and $x(t)$ are $m \times 1$ and $n \times 1$ vectors of source and recorded signals, respectively, and $A \in R^{n \times m}$ is a mixing matrix.
- The goal of blind source separation (BSS) is then to find $s(t)$ only by observing $x(t)$. The general BSS problem is impossible for the case $n<m$. However, if the sources are sparse then this problem can be solved (using L1 minimization).
- We have the problem of finding sparse solutions of the undetermined system of linear equations (USLE) $A s=x$. To obtain the sparsest solution of $A s=x$, we may search for a solution with minimal L0 norm. (Intractable problem, sensitive to noise)
- Hence, researchers consider L1 approaches such as basis pursuit (BP), LP-norm approaches such as IRLS, and greedy approaches such as matching pursuit (MP).
- In this paper, authors present an approach for solving USLE by direct minimization of the L0 norm after smoothing (approximating with smooth functions).
- Performance of the algorithm is equal to (or better than) the interior point based algorithms with 2 to 3 orders of magnitude faster.

| algorithm | total time $(\mathrm{sec})$ | MSE |
| :---: | :---: | :---: |
| SL0 | 0.227 | $5.53 e-5$ |
| LP $\left(\ell_{1}\right.$-magic $)$ | 30.1 | $2.31 e-4$ |
| FOCUSS | 20.6 | $6.45 e-4$ |

## II. Approach

- L0 norm of a vector $x \in R^{m}$ is a discontinuous function of that vector.

$$
\|x\|_{0}=\sum_{i=1}^{m} I_{i} \quad I_{i}= \begin{cases}1 & \text { if } x_{i} \neq 0 \\ 0 & \text { if } x_{i}=0\end{cases}
$$

- The idea then is to approximate the discontinuous function with a continuous function. The continuous function has a parameter ( $\operatorname{say} \sigma$ ) that determines the quality of the approximation.
- For example, consider the (one-variable) family of Gaussian functions

$$
f_{\sigma}(s) \triangleq \exp \left(\frac{-s^{2}}{2 \sigma^{2}}\right)
$$

and note that

$$
\lim _{\sigma \rightarrow 0} f_{\sigma}(s)= \begin{cases}1 & \text { if } s=0 \\ 0 & \text { if } s \neq 0\end{cases}
$$

or approximately

$$
f_{\sigma}(s) \approx \begin{cases}1 & \text { if }|s| \ll \sigma \\ 0 & \text { if }|s| \gg \sigma\end{cases}
$$

- Now define

$$
\begin{aligned}
& F_{\sigma}(s)=\sum_{i=1}^{m} f_{\sigma}\left(s_{i}\right) \\
& \|s\|_{0} \approx m-F_{\sigma}(s)
\end{aligned}
$$

- For small values of $\sigma$, the approximation tends to equality. Hence, we can define the minimum L0 norm solution by maximizing $F_{\sigma}(s)$.
- The value of $\sigma$ determines how smooth the function $F_{\sigma}(s)$ is: the larger value of $\sigma$, the smoother $F_{\sigma}(s)$ (but worse approximation to L0-norm); and the smaller value of $\sigma$, closer the behavior of $F_{\sigma}(s)$ to L0-norm.
- However, for smaller values of $\sigma, F_{\sigma}(s)$ is highly non-smooth and contains a lot of local maxima, and hence its maximization is not easy. On the other hand, for larger values of $\sigma, F_{\sigma}(s)$ is smoother and contains less local maxima (in fact, no local maxima for large $\sigma$ ).
- "Basic idea": In order to find an $s$ that maximizes $F_{\sigma}(s)$, the authors start with maximum $\sigma$. For this maximum $\sigma$, they find the maximizer of $F_{\sigma}(s)$. Then they decrease $\sigma$ and again find the maximizer of $F_{\sigma}(s)$.
- They claim that eventually this process (decreasing $\sigma$ and maximizing $F_{\sigma}(s)$ ) results in the maximization of $F_{\sigma}(s)$ or equivalently minimization of the L 0 norm.
- Other family of functions that approximates the Kronecker delta functions like family of triangular functions,

$$
f_{\sigma}(s)= \begin{cases}1, & \text { if }|s| \geq \sigma \\ \frac{\sigma \sigma \sigma s}{}, & \text { if }-\sigma \leq s \leq 0 \\ \frac{(\sigma-a s)}{\sigma}, & \text { if } 0 \leq s \leq \sigma\end{cases}
$$

- or truncated hyperbolic functions

$$
f_{\sigma}(s)= \begin{cases}1, & \text { if }|s| \geq \sigma \\ 1-\left(\frac{s}{\sigma}\right)^{2}, & \text { if }|s| \leq \sigma\end{cases}
$$

- or functions of the form

$$
f_{\sigma}(s)=\frac{\sigma^{2}}{\left(s^{2}+\sigma^{2}\right)} .
$$

- For sufficiently large values of $\sigma$ the maximizer of $F_{\sigma}(s)$ subject to $A s=x$ is the minimum L2-norm solution, i.e., $\hat{s}=A^{T}\left(A A^{T}\right)^{-1} x$. Justification of the statement for Gaussian family:

We want to maximize $F_{\sigma}(s)=\sum_{i=1}^{m} f_{\sigma}\left(s_{i}\right)=\sum_{i=1}^{m} e^{\left(-s_{i}^{2} / 2 \sigma^{2}\right)}$ subject to $A s=x$
The Lagrangian is $F_{\sigma}(s)-\lambda^{T}(A s-x)$. Differentiating the Lagrangian w.r.t $s$ and $\lambda$ and setting the result to zero gives the following KKT systems of $m+n$ non-linear equations of $m+n$ unknowns.

$$
\begin{aligned}
& {\left[s_{1} e^{-x_{1}^{2} / 2 \sigma^{2}}, s_{2} e^{-s^{2} / 2 \sigma^{2}}, \cdots, s_{m} e^{-\frac{z^{2} / 2 \sigma^{2}}{2}}\right]^{T}-A^{T} \lambda_{1}=0 \quad \lambda_{1}=-\sigma^{2} \lambda} \\
& A s-x=0
\end{aligned}
$$

Now, let us look at this problem: $\min \|s\|_{2}^{2}$ s.t. $A s=x$. Again using Lagrange multipliers this minimization results in the system of equations

$$
\begin{aligned}
& {\left[s_{1}, s_{2}, \cdots, s_{m}\right]^{T}-A^{T} \lambda=0} \\
& A s-x=0
\end{aligned}
$$

Authors claim that for a larger $\sigma$ these two systems of equations are identical and hence the maximizer of $F_{\sigma}(s)$ for larger $\sigma$ is the minimum L2 norm solution.

- Hence, the authors start with large $\sigma$ and maximize the corresponding $F_{\sigma}(s)$. They then decrease $\sigma$ and repeat the maximization of $F_{\sigma}(s)$ again. They repeat the process for a few sequences of $\sigma$ and shown that the subsequent maximization of $F_{\sigma}(s)$ leads to L0 solution. Their algorithm is based on the principles of graduated non-convexity.

From Wikipedia:

- Graduated non-convexity is a global optimization technique that attempts to solve a difficult optimization problem by initially solving a greatly simplified problem, and progressively transforming that problem (while optimizing) until it is equivalent to the difficult optimization problem.
- Graduated optimization is an improvement to hill climbing that enables a hill climber to avoid settling into local optima. It breaks a difficult optimization problem into a sequence of optimization problems, such that the first problem in the sequence is convex (or nearly convex), the solution to each problem gives a good starting point to the next problem in the sequence, and the last problem in the sequence is the difficult optimization problem that it ultimately seeks to solve. Often, graduated optimization gives better results than simple hill climbing [1].



## III. The SL0 Algorithm

The algorithm consists of 2 loops: the outer and inner loops.
In the outer loop, they vary the values of $\sigma$.
In the inner loop, for a given $\sigma$ they use a steepest ascent algorithm for maximizing $F_{\sigma}(s)$.

## - Initialization:

1) Let $\mathrm{s}_{0}$ be equal to the minimum $\ell^{2}$ norm solution of $\mathbf{A s}=\mathrm{x}$, obtained by pseudo-inverse of $\mathbf{A}$.
2) Choose a suitable decreasing sequence for $\sigma,\left[\sigma_{1} \ldots \sigma_{J}\right]$ (see Remarks 5 and 6 of the text).

- For $j=1, \ldots, J$ :

1) Let $\sigma=\sigma_{j}$.
2) Maximize (approximately) the function $F_{\sigma}$ on the feasible set $\mathcal{S}=\{\mathbf{s} \mid \mathbf{A s}=\mathbf{x}\}$ using $L$ iterations of the steepest ascent algorithm (followed by projection onto the feasible set):

- Initialization: $\mathbf{s}=\hat{\mathbf{s}}_{j-1}$.
- For $\ell=1 \ldots L$ (loop $L$ times):
a) Let $\delta \triangleq\left[s_{1} \exp \left(-s_{1}^{2} / 2 \sigma^{2}\right), \ldots, s_{n} \exp \left(-s_{n}^{2} / 2 \sigma^{2}\right)\right]^{T}$.
b) Let $\mathrm{s} \leftarrow \mathrm{s}-\mu \delta$ (where $\mu$ is a small positive constant).
c) Project s back onto the feasible set $\mathcal{S}$ :

$$
\mathbf{s} \leftarrow \mathbf{s}-\mathbf{A}^{T}\left(\mathbf{A} \mathbf{A}^{T}\right)^{-1}(\mathbf{A} \mathbf{s}-\mathbf{x})
$$

3) Set $\hat{\mathbf{s}}_{j}=\mathbf{s}$.

- Final answer is $\hat{\mathbf{s}}=\hat{\mathbf{s}}_{J}$.

Fig. 1. Final SL0 algorithm.
Remark 1: The internal loop is repeated for a fixed, say $L$, and small number of times. In other words, we do not wait for the steepest ascent algorithm (SAA) to converge. That is, we do not need the exact maximizer of $F_{\sigma}(s)$. We just need to enter the region near the (global) maximizer of for escaping from its local maximizer.

Remark 2: The SAA consists of update step $s \leftarrow s+\mu_{j} \nabla F_{\sigma}(s)$. Here, $\mu_{j} \mathrm{~s}$ are step size parameter and they should be chosen such that for decreasing values of $\sigma, \mu_{j}$ should be smaller. In the algorithm, they let $\mu_{j}=\mu \sigma^{2}$ for some constant $\mu$. Then, $s \leftarrow s+\mu_{j} \nabla F_{\sigma}(s)=s-\mu \delta$ where $\delta=-\sigma^{2} \nabla F_{\sigma}=\left[s_{1} e^{-s_{1}^{2} / 2 \sigma^{2}}, s_{2} e^{-s_{2}^{2} / 2 \sigma^{2}}, \cdots, s_{m} e^{-s_{m}^{2} / 2 \sigma^{2}}\right]^{T}$.

Remark 3: Each iteration of the inner loop consists of gradient ascent step, followed by a projection step.
If we are looking for a suitable large $\mu$ (to reduce the required number of iterations), a suitable choice is to make the algorithm to force all those values of $s_{i}$ satisfying $\left|s_{i}\right| \leq \sigma$ toward zero. For this aim, we should have $\mu \exp \left(\frac{-s_{i}^{2}}{2 \sigma^{2}}\right) \approx 1$, and because $\exp \left(\frac{-s_{i}^{2}}{2 \sigma^{2}}\right) \leq 1$ for $\left|s_{i}\right| \leq \sigma$, the choice $\mu \geq 1$ seems reasonable.

Remark 4: The algorithm may work by initializing to an arbitrary solution. However, as discussion before the best initial value of is the minimum L2 solution. In fact, calculating minimum L2norm is one of the earliest approaches used for estimating the sparsest solution called the method of frames [4].

Remark 5: After initiating with minimum L2 solution, the next value for $\sigma$ may be chosen about two to four times of the maximum absolute value of the obtained solution $\left(\max _{i}\left|s_{i}\right|\right)$. For example, if we take $\sigma>4 \max _{i}\left|s_{i}\right|$, then $\exp \left(\frac{-s_{i}^{2}}{2 \sigma^{2}}\right) \geq 0.96 \approx 1 i=1,2, \cdots, m$. This value of $\sigma$ acts virtually like infinity for all values of $s_{i}$.

Remark 6: The term $F_{\sigma}(s)\left(\|s\|_{0} \approx m-F_{\sigma}(s)\right)$ simply counts the number of zero components of $s$. However, instead of hard-thresholding that is "zero $\equiv\left|s_{i}\right|<\sigma$ " and "non-zero $\equiv\left|s_{i}\right|>\sigma " \quad f_{\sigma}(s) \approx\left\{\begin{array}{ll}1 & \text { if }|s| \ll \sigma \\ 0 & \text { if }|s| \gg \sigma\end{array}\right.$ uses a soft-thresholding, for which $\sigma$ is a rough threshold.

## Remark 7:

- If $\boldsymbol{s}$ is an exactly $K$-sparse signal, then $\sigma$ can be decreased to arbitrarily small values. In fact, in this case, the minimum value of $\sigma$ is determined by the desired accuracy (as will be discussed in Theorem 1).
- If $s$ is an approximately $K$-sparse signal (say the source vector is noisy), then the smallest $\sigma$ should be about one to two times of (a rough estimation of) the standard deviation of the noise (in the source vector). This is because, while $\sigma$ is in this range, $f_{\sigma}(s)$ shows that the cost function treats small (noisy) samples as zeros (i.e., for which $f_{\sigma}(s) \approx 1$ ).
- However, below this range, the algorithm tries to 'learn' these noisy values, and moves away from the true answer. (According to the previous remark, the soft threshold should be such that all these noisy samples be considered zero).
- Restricting $\sigma$ to be above the standard deviation of the noise, provides the robustness of this approach to noisy sources, which was one of the difficulties in using the exact L0 norm.


## IV. ANALYSIS OF THE ALGORITHM

## A. Convergence analysis

In this section, we try to answer two questions for the noiseless case (the noisy case will be considered in Section IV-C):
a) Does the SL0 solution converges to the actual minimizer of the L0 norm?
b) If yes, how much should we decrease $\sigma$ to achieve a desired accuracy?

Assuming the maximization of $F_{\sigma}(s)$ for fixed $\sigma$ is perfectly done (and we obtain the maximizer $s^{\sigma}$ ). The authors show that the sequence of 'global' maximizers of $F_{\sigma}(s)$ 's will converge to the sparsest solution. That is we need to prove

$$
\lim _{\sigma \rightarrow 0} s^{\sigma}=s^{0}
$$

For proving the above statement the authors have introduced three intermediate results via lemmas.

Lemma 1: Assume a matrix $A \in R^{n \times m}$ has the property that all of its $n \times n$ sub-matrices are invertible, which is called the unique representation property (URP) in [3]. For any $\forall s \in N(A)$ if the $m-n$ elements of $s$ have absolute values less than $\alpha \Rightarrow\|s\| \leq \beta$.

Proof: We have to show that
$\forall \beta>0, \quad \exists \alpha>0, \quad$ s.t $\forall s \in N(A):$
$m-n$ elements of $s$ have absolute values less than $\alpha \Rightarrow\|s\| \leq \beta \quad\|\cdot\|$ stands for L 2 norm
Let $s \in N(A)$ and assume that the absolute values of at least $m-n$ elements of it are smaller than $\alpha$. Let $I_{\alpha}$ be the set of all indices $i$, for which $\left|s_{i}\right|>\alpha$.Consequently, $\left|I_{\alpha}\right| \leq n$. Then we write

$$
\begin{aligned}
& \sum_{i=1}^{m} s_{i} a_{i}=0 \Rightarrow \sum_{i \in I_{\alpha}} s_{i} a_{i}+\sum_{i \neq I_{\alpha}} s_{i} a_{i}=0 \\
& \Rightarrow\left\|\sum_{i \in I_{\alpha}} s_{i} a_{i}\right\|=\left\|\sum_{i \notin I_{\alpha}} s_{i} a_{i}\right\| \leq \sum_{i \notin I_{\alpha}}\left\|s_{i} a_{i}\right\| \\
& =\sum_{i \neq I_{\alpha} \leq \alpha=1} s_{i} \mid\left\|a_{i}\right\| \\
& \leq\left(m-\left|I_{\alpha}\right|\right) \alpha \leq m \alpha
\end{aligned}
$$

Let $\hat{\mathrm{A}}$ be the submatrix of $A$ containing only those columns of $A$ that are indexed by the elements of $I_{\alpha}$. Thus, A has at most $n$ columns, and the columns of $\hat{\mathrm{A}}$ are linearly independent, because of the URP of $A$. Therefore, there exists a left inverse $\hat{\mathrm{A}}^{-1}$ for $\hat{\mathrm{A}}$. Let $\bar{s}$ and $\tilde{s}$ denote those sub-vectors of $\boldsymbol{s}$ which are, and which are not indexed by $I_{\alpha}$, respectively. Then

$$
\begin{aligned}
& \left(\sum_{i \in I_{\alpha}} s_{i} a_{i}\right)=\hat{\mathrm{A}} \overline{\mathrm{~s}} \Rightarrow \overline{\mathrm{~s}}=\hat{\mathrm{A}}^{-1}\left(\sum_{i \in I_{\alpha}} s_{i} a_{i}\right) \\
& \|\overline{\mathrm{s}}\|=\left\|\hat{\mathrm{A}}^{-1}\left(\sum_{i \in I_{\alpha}} s_{i} a_{i}\right)\right\| \\
& \|\overline{\mathrm{s}}\| \leq\left\|\hat{\mathrm{A}}^{-1}\right\|\left\|\sum_{i \in I_{\alpha}} s_{i} a_{i}\right\|=\left\|\hat{\mathrm{A}}^{-1}\right\| m \alpha \\
& \|\tilde{\mathrm{~s}}\| \leq \sum_{i \neq I_{\alpha}}\left|s_{i}\right| \leq\left(m-\left|I_{\alpha}\right|\right) \alpha \leq m \alpha \\
& \|\mathrm{~s}\| \leq\|\overline{\mathrm{s}}\|+\|\tilde{\mathrm{s}}\| \leq\left(1+\left\|\hat{\mathrm{A}}^{-1}\right\|\right) m \alpha
\end{aligned}
$$

Now, let M be the set of all submatrices $\hat{\mathrm{A}}$ of $A$, consisting of at most $n$ columns of $A$. Then, $M$ is clearly a finite set (in fact $|M|<2^{m}$ ).
Let $M=\max \left\{\left\|\hat{A}^{-1}\right\| \mid \hat{A} \in M\right\}$ then

$$
\|\mathrm{s}\| \leq\left(1+\left\|\hat{\mathrm{A}}^{-1}\right\|\right) m \alpha \leq(1+M) m \alpha
$$

$M$ is a constant and its value depends only on the matrix $A$. Thus, for each $\beta$ it is suffice to choose $\alpha=\beta / m(M+1)$

Corollary 1: If $A \in R^{n \times m}$ satisfies the URP, and $s \in N(A)$ has at most $n$ elements with absolute values greater than $\alpha$, then $\|\mathrm{s}\|<(1+M) m \alpha$.

## Lemma 2:

Let a function $f_{\sigma}(s)$ have the properties $f_{\sigma}(0)=1$ and $\forall s 0 \leq f_{\sigma}(s) \leq 1$, and let $F_{\sigma}(s)=\sum_{i=1}^{m} f_{\sigma}\left(s_{i}\right)$. Assume $A$ satisfies the URP, and let $S=\{s \mid A s=x\}$. Assume that there exists a (sparse) solution $s^{0} \in S$ for which $\|s\|_{0}=k \leq \frac{n}{2}$ (such a sparse solution is unique). Then, if for a solution $\hat{s}=\left(\hat{s}_{1}, \hat{s}_{2}, \cdots, \hat{s}_{m}\right)^{T} \in S, F_{\sigma}(\hat{s}) \geq m-(n-k)$ and if $\alpha>0$ is chosen such that the $\hat{s}_{i}{ }^{\prime} s$ with absolute values greater than $\alpha$ satisfy $f_{\sigma}\left(\hat{s}_{i}\right) \leq 1 / m$, then $\left\|\hat{s}-s^{0}\right\|<(M+1) m \alpha$

Proof: Let $I_{\alpha}$ be the set of all indices $i$, for which $\left|\hat{s}_{i}\right|>\alpha$, then

$$
\begin{aligned}
F_{\alpha}(\hat{s}) & =\sum_{i=1}^{m} f_{\alpha}\left(\hat{s}_{i}\right) \\
& =\underbrace{\left.f_{i}\right)}_{i_{i I_{\alpha}} \underbrace{f_{\alpha}\left(\hat{s_{i}}\right)}_{\leq m .(1 / m)=1}}+\underbrace{\sum_{i \neq I_{\alpha}} \underbrace{f_{\alpha}\left(\hat{s}_{i}\right)}_{\leq 1}}_{\leq m-\left|I_{\alpha}\right|} \\
& \leq 1+m-\left|I_{\alpha}\right|
\end{aligned}
$$

We assume that we have chosen $f_{\sigma}(s)$ such that $F_{\sigma}(\hat{s}) \geq m-(n-k)$. (We prove this next) Now, we get $m-(n-k) \leq F_{\sigma}(\hat{s}) \leq 1+m-\left|I_{\alpha}\right|$, from which we can get $\left|I_{\alpha}\right| \leq n-k$.

As a result, at most $n-k$ elements of $\hat{s}$ have absolute values greater than $\alpha$.
Since $s^{0}$ has exactly $k$ non-zero elements, we conclude that $\hat{s}-s^{0}$ has at most $(n-k)+k=n$ elements with absolute values greater than $\alpha$.

Moreover, $\hat{s}-s^{0} \in N(A)$ and hence by Corollary 1 we have $\left\|\hat{s}-s^{0}\right\|<(M+1) m \alpha$.
Corollary 2: For the Gaussian family $f_{\sigma}(s) \triangleq \exp \left(\frac{-s^{2}}{2 \sigma^{2}}\right)$, if $F_{\sigma}(\hat{s}) \geq m-(n-k)$ holds for a solution $\hat{s}$, then

$$
\left\|\hat{s}-s^{0}\right\|<(M+1) m \sigma \sqrt{2 \ln m}
$$

## Proof:

For the Gaussian family $f_{\sigma}(s) \triangleq \exp \left(\frac{-s^{2}}{2 \sigma^{2}}\right), \alpha$ required for lemma 2 can be chosen as $\alpha=\sigma \sqrt{2 \ln m}$. Because for $\left|\hat{s}_{i}\right|>\sigma \sqrt{2 \ln m}$,

$$
f_{\sigma}\left(\hat{s}_{i}\right)=\exp \left(\frac{-\hat{s}_{i}^{2}}{2 \sigma^{2}}\right)<\exp \left(\frac{-2 \sigma^{2} \ln m}{2 \sigma^{2}}\right)=\frac{1}{m}
$$

Moreover, Gaussian family satisfies the other condition required in lemma 2.
Lemma 3: Let $f_{\sigma}, F_{\sigma}, S$ and $s^{0}$ be as in Lemma 2, and let $s^{\sigma}$ be the maximizer of $F_{\sigma}(s)$ on $S$, then $s^{\sigma}$ satisfies $F_{\sigma}(\hat{s}) \geq m-(n-k)$.
Proof: We write

$$
\begin{aligned}
F_{\sigma}\left(s^{\sigma}\right) & \geq F_{\sigma}\left(s^{0}\right) \\
& \geq m-k \\
& \geq m-(n-k) \quad\left(\because k \leq \frac{n}{2}\right)
\end{aligned}
$$

Note that Lemma 3 and Corollary 2 prove together that for the Gaussian family $f_{\sigma}(s) \triangleq \exp \left(\frac{-s^{2}}{2 \sigma^{2}}\right) \underset{A s=x}{\arg \max } F_{\sigma}(s) \rightarrow s^{0}$ as $\sigma \rightarrow 0$. This result can however be stated for a larger class of functions, as done in Theorem 1 (next page).

Theorem 1: Consider a family of univariate functions $f_{\sigma}$, indexed $\sigma \in R^{+}$satisfying the following set of conditions:

1. $\lim _{\sigma \rightarrow 0} f_{\sigma}(s)=0 ; \quad \forall s \neq 0$
2. $f_{\sigma}(0)=1 ; \quad \forall \sigma \in R^{+}$
3. $0 \leq f_{\sigma}(s) \leq 1 ; \quad \forall \sigma \in R^{+} \quad s \in R$
4. For each positive values of $v$ and $\alpha$, there exists $\sigma_{0} \in R^{+}$that satisfies

$$
\begin{equation*}
|s|>\alpha \Rightarrow \quad f_{\sigma}(s)<v ; \quad \forall \sigma<\sigma_{0} \tag{1}
\end{equation*}
$$

Assume $A$ satisfies the URP, and let $F_{\sigma}, S$ and $s^{0}$ be defined as in Lemma 2, and $s^{\sigma}=\left(s_{1}^{\sigma}, s_{2}^{\sigma}, \cdots, s_{m}^{\sigma}\right)^{T} \in S$ be the maximizer of $F_{\sigma}$ on $S$. Then:

$$
\lim _{\sigma \rightarrow 0} s^{\sigma}=s^{0}
$$

## Working definition of limit of a sequence

We say that $\lim _{n \rightarrow \infty} a_{n}=L$ if we can make $a_{n}$ as close to $L$ as we want for sufficiently large $n$.

## Precise definition of the limit

We say that $\lim _{n \rightarrow \infty} a_{n}=L$ if for some positive error term $\varepsilon$ the distance of the sequence at $n$ from $L$ must be less than the allowed error $\varepsilon$, that is, $\left|a_{n}-L\right|<\varepsilon$. But, it is important to remember that it is not enough that our sequence does converge once or twice; it must be within the error for all values from some point onwards, that is, $\left|a_{n}-L\right|<\varepsilon, \forall n>N$.

Analytically, $\quad \forall \varepsilon>0, \quad \exists N \in \mathbb{N}, \quad \forall n>N \quad\left|a_{n}-L\right|<\varepsilon$.

Proof: To prove $\lim _{\sigma \rightarrow 0} s^{\sigma}=s^{0}$, we have to show that

$$
\begin{equation*}
\forall \beta>0, \quad \exists \sigma_{0}>0, \quad \forall \sigma<\sigma_{0} \quad\left\|s^{\sigma}-s^{0}\right\|<\beta \tag{2}
\end{equation*}
$$

For each $\beta>0$, let $\alpha=\beta / m(M+1)$. Then for this $\alpha$ and $\nu=1 / m$, condition 4 of the theorem gives a $\sigma_{0}$ for which the (1) holds. We show that this is the $\sigma_{0}$ we were seeking for in (2).

Note that $\forall \sigma<\sigma_{0}$, (1) states that for $s_{i}^{\sigma}$ 's with absolute values greater than $\alpha$ we have $f_{\sigma}\left(s_{i}^{\sigma}\right) \leq 1 / m$. Moreover, Lemma 3 states that $s^{\sigma}$ satisfies $F_{\sigma}\left(s^{\sigma}\right) \geq m-(n-k)$. Consequently, all the conditions of Lemma 2 have been satisfied, and hence it implies that $\left\|s^{\sigma}-s^{0}\right\|<(M+1) m \alpha=\beta$.

Remark 1: The Gaussian family $f_{\sigma}(s) \triangleq \exp \left(\frac{-s^{2}}{2 \sigma^{2}}\right)$ satisfies conditions 1 through 4 of Theorem 1. Other Families of functions also satisfy the conditions of Theorem 1.

Remark 2: Using Corollary 2, where using Gaussian family, to ensure an arbitrary accuracy in estimation of the sparse solution $s^{0}$, it suffices to choose

$$
\sigma<\frac{\beta}{m(M+1) \sqrt{2 \ln m}} \quad \text { and do the optimization of } F_{\sigma}(s) \text { subject to } A s=\boldsymbol{x} .
$$

Remark 3: Consider the set of solutions in $\hat{s}^{\sigma}$ in $S$, which might not be the absolute maxima of functions $F_{\sigma}$ on $S$, but satisfy the condition

$$
F_{\sigma}\left(\hat{s}^{\sigma}\right) \geq m-(n-k)
$$

By following a similar approach to the proof of Theorem 1, it can be proved that $\lim _{\sigma \rightarrow 0} s^{\sigma}=s^{0}$. In other words, for the steepest ascent (internal loop), it is not necessary to reach the absolute maximum. It is enough reach a solution in which is $F_{\sigma}$ large.

Remark 4: The previous remark proposes another version of SL0 in which there is no need to set a parameter L: Repeat the internal loop until $F_{\sigma}(s)$ exceeds $m-n / 2$ [the worst case of the limit given by $\left.F_{\sigma}(\hat{s}) \geq m-(n-k)\right]$ or $m-(n-k)$ if $k$ is known a priori. The advantage of such a version is that if it converges, then it is guaranteed that the estimation error is bounded as $\left\|\hat{s}-s^{0}\right\|<(M+1) m \sigma \sqrt{2 \ln m}$, in which $\sigma$ is replaced with $\sigma_{J}$, the last element of the sequence of $\sigma$.

It has, however, two disadvantages: first, it slows down the algorithm because exceeding the limit $m-(n-k)$ for each $\sigma$ is not necessary (it is just sufficient); and second, because of the possibility that the algorithm runs into an infinite loop because $F_{\sigma}(s)$ cannot exceed this limit (this occurs if the chosen sequence of $\sigma$ has not been resulted in escaping from local maxima).

Remark 5: As another consequence, Lemma 1 provides an upper bound on the estimation error $\left\|\hat{s}-s^{0}\right\|$, only by having an estimation $\hat{s}$ (which satisfies $A \hat{s}=x$ ): Begin by sorting the elements of $\hat{s}$ in descending order and $\alpha$ let be the absolute value of the $\left\lfloor\frac{n}{2}\right\rfloor+1$ 'th element. Since $s^{0}$ has at most $n / 2$ nonzero elements, $\hat{s}-s^{0}$ has at most $n$ elements with absolute values greater than $\alpha$. Moreover, $\hat{s}-s^{0} \in N(A)$ and hence Corollary 1 implies that $\left\|\hat{s}-s^{0}\right\|<(M+1) m \alpha$. This result is consistent with the heuristic that "if $\hat{s}$ has at most $n / 2$ 'large' components, the uniqueness of the sparsest solution insures that $\hat{s}$ is close to the true solution."

## B. Relation to minimum norm 2 solution

Kindly refer the paper for proof of $\lim _{\sigma \rightarrow \infty} \underset{A s=x}{\arg \max } F_{\sigma}(s)=\hat{s}$, where $\hat{s}$ is the minimum L2 norm solution.
C. Noisy case

As shown in the proof of Theorem 1 (noiseless case), a smaller value of $\sigma$ results in a more accurate solution and it is possible to achieve solutions as accurate as desired by choosing small enough values of $\sigma$. However, this is not the case in the presence of additive noise, that is, if $x=\mathrm{A} s+n$. In fact, the noise power bounds the maximum achievable accuracy.

## V. Numerical Results

The performance of the SL0 algorithm is experimentally verified and is compared with BP (FOCUSS) and LP (L1 magic). The effects of the parameters, sparsity, noise, and dimension on the performance are also experimentally discussed (Please refer the paper).

In experiments, sparse sources are artificially created using a Bernoulli-Gaussian model: each source is "active" with probability $p$, and is "inactive" with probability ( $1-p$ ). If it is active, each sample is a zero-mean Gaussian random variable with variance $\sigma_{\text {on }}^{2}$; if it is not active, each sample is a zero-mean Gaussian random variable with variance $\sigma_{\text {off }}^{2}$, where $\sigma_{\text {off }}^{2} \ll \sigma_{\text {on }}^{2}$.

Each column of the mixing matrix is randomly generated using the normal distribution and then is normalized to unity.

To evaluate the estimation quality, signal-to-noise ratio (SNR) and mean-square error (MSE) are used. SNR (in dB) is defined as $20 \log \left(\frac{\|s\|}{\|s-\hat{s}\|}\right)$ and MSE as $\frac{1}{m}\|s-\hat{s}\|^{2}$

TABLE I
Progress of SL0 For a Problem With $m=1000, n=400 \mathrm{AND}$
$k=100(p=0.1)$

| itr. \# | $\sigma$ | MSE | SNR (dB) |
| :---: | :---: | :---: | :---: |
| 1 | 1 | $4.84 e-2$ | 2.82 |
| 2 | 0.5 | $2.02 e-2$ | 5.19 |
| 3 | 0.2 | $4.96 e-3$ | 11.59 |
| 4 | 0.1 | $2.30 e-3$ | 16.44 |
| 5 | 0.05 | $5.83 e-4$ | 20.69 |
| 6 | 0.02 | $1.17 e-4$ | 28.62 |
| 7 | 0.01 | $5.53 e-5$ | 30.85 |
| algorithm | total time (sec) | MSE | SNR (dB) |
| SL0 | 0.227 | $5.53 e-5$ | 30.85 |
| LP ( $\ell_{1}-$ magic $)$ | 30.1 | $2.31 e-4$ | 25.65 |
| FOCUSS | 20.6 | $6.45 e-4$ | 20.93 |

## VI. Conclusions

In this paper, authors showed that the smoothed L0 norm can be used for finding sparse solutions of an USLE. They also showed that the smoothed version of the L0 norm results in an algorithm which is faster than the state-of-the-art algorithms based on minimizing the L1 norm.

Moreover, this smoothing solves the problem of high sensitivity of L0 norm to noise. In another point of view, the smoothed L0 provides a smooth measure of sparsity.

The basic idea of the paper was justified by both theoretical (convergence in both noiseless and noisy case, relation to the L2 norm solution) and experimental analysis of the algorithm.

## Appendix (Not available in the paper)

The following are taken from [5]
Consider the following problem of the Euclidean orthogonal projection of a point to an affine set: For the given $A \in R^{m \times n}, b \in R^{m}$ and $p \in R^{n}$, find a vector $x^{*} \in R^{n}$ satisfying

$$
\begin{aligned}
& A x^{*}=b \\
& \left\|p-x^{*}\right\|=\min _{A x=b}\|p-x\|
\end{aligned}
$$

The solution to the above problem exits and it is unique and it is

$$
\begin{array}{rlr}
x^{*} & =p-A^{+} A p+A^{+} b & {\left[x^{*}=p-A^{+}(A p-b)\right]} \\
& =\left[I-A^{+} A\right] p+A^{+} b & \\
x^{*} & =P_{N(A)} p+A^{+} b &
\end{array}
$$

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